

## Universidad de Valladolid

# **Support for Thread-Level Speculation into** OpenMP

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# INTRODUCTION

- Manual development of parallel versions of sequential applications is a **difficult task**. It requires:
  - In-depth knowledge of the problem.
  - Understanding of the underlying architecture.
  - Knowledge on the parallel programming model.
- **OpenMP** allows to parallelize code "avoiding" these requirements.
- Compilers' **automatic parallelization** only proceed when there is no risk.
- Thread-Level Speculation (TLS) can extract parallelism when a compile-time dependence analysis can not guarantee that the code is safely parallelizable.
- We have already developed a TLS runtime library.
- Current goal: To automatically transforms loops written in OpenMP syntax to benefit from speculative parallelization.

### Our proposal

- Goal: Add TLS support into OpenMP.
- New OpenMP clause:

#pragma omp parallel for ` speculative(variable/,var\_list/)

- **Speculative variables** are those whose use can potentially lead to a dependence violation. They need to be monitored at runtime in order to obtain results.
- Programmer classifies variables in private, shared, and a new category: **speculative**.
- TLS should be **transparent** from the point of view of the programmer. If he/she is unsure about the use of a certain structure, he/she could simply label it as speculative. The compiler automatically will transform the code in order to speculatively parallelize the loop.

# HOW THREAD-LEVEL SPECULATION WORKS?

Thread 1 (non spec)	Thread 2	Thread 3	Thread 4 (most-spec)
(iter. 1, x = 1)	(iter. 2, x = 1)	(iter. 3, x = 2)	(iter. 4, x = 2)



Our TLS is implemented using OpenMP for thread management.

### ARCHITECTURE

### MODIFYING GCC

- Use GCC as reference compiler.
- Since version 4.5, GCC can be extended by **plugins**:
  - Faster prototyping.
  - Easier modifications.
  - Extensibility: new compiler passes.
- The parser recognizes the new clause, and a new pass performs the transformations needed.
- Transformations are done before the compiler optimization passes.
- The new pass works with the **GIMPLE** representation.

#### Middle End Back End Front End C parser GIMPLE RTL C++ parser RTL Interprocedural Optimizer Optimizer GENERIC

#### Fortran parser Final Code SSA Manager Generation Optimizer **Plugin pass** Assembly Java parser \* main OpenMP related parts

## CONCLUSIONS

• Adding speculative support to OpenMP would greatly increase the number of loops that could be parallelized with this programming model.

Call Graph

Manager

Pass

• The programmer may label some of the variables involved as private or shared, using spec-

**Original annotated** 

### **Code generated**

- $\rightarrow$  specinit();
- -> omp\_set\_num\_threads(T);
- > specstart(N);
  - #pragma omp parallel for \ private(a), shared (b) \

CODE EXAMPLE

**#pragma omp parallel for** \

private(a), shared (b) \

 $-\langle - \rangle$  private(engine\_vars), shared(engine\_vars) \ speculative(v) - - -→ shared(v)

initSpecLoop(v, 1); **for** (i=0; i < N; i++) **{** specload(a, v, i); a = v[i]; specstore(v, i, b); v[i] = b; endSpecLoop(v, N);

• This transformations are done by the new pass **automatically**.

- It detects each reading from and writing into the speculative variable and replaces them for specstore() and specload() functions.
- It also add all the structures and functions needed to speculatively parallelize the code.

#### Acknowledgements

ulative for the rest.

- The parser detects the new speculative clause, and the new compiler pass **performs** automatically all the transformations needed to speculatively parallelize the loop.
- This process is **transparent to program**mers. They do not need to know anything about the speculative parallelizing model. • Our proposal would let to transform *any* loop into a parallel loop.

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